Graphs for Genomic Sequences

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At the beginning ...

... there was the quest for understanding THE GENOME:

- reading the genome sequence (composed of 4 nucleotides A, T, G, C) followed by
- understanding the functions of each piece
- classifying the variations among genomes

HOW: by reading genomes from different species and individuals and comparing them

PROBLEM: this is extremely difficult



culminated with the Human Genome Project (HGP) officially started in 1990

The human genome

"is a book with multiple uses; a history book - a narrative of the journey of our species through time. It's a shop manual, with an incredibly detailed blueprint for building every human cell. And it's a transformative textbook of medicine, with insights that will give health care providers immense new powers to treat, prevent and cure disease"

Francis Collins, director of the "National Human Genome Research Institute"

The human genome

The HGP published the first draft of the human genome in Nature, Feb. 2001

- 3 billion base pairs (90% percent complete)
- detailed information about the structure, organization and function of the complete set of human genes
- about 20,500 genes (significantly fewer than the previous estimates ranging from 50,000 to 140,000 genes)
- the full sequence was completed and published in April 2003.

The human genome



• the full sequence was completed and published in April 2003.

20 years: from HGP to 1000 Genomes Project

A global reference for human genetic variation

The 1000 Genomes Project Consortium

Nature volume 526, pages 68–74 (01 October 2015)

1000 Genomes Project

A global reference for human genetic variation

The 1000 Genomes Project Consortium* OCTOBER 2015 | VOL 526 | NATURE

Phase III: 2,504 humans : 84.8 million SNPs





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Different people ...



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The human genome book is written in a 4 - letter alphabet, books differing slightly between different people :

- roughly 1 difference in 1000 nucleotides in average, accounting from height to genetic diseases a typical genome differs from the "reference human genome" in ~ 5 million places, ~ 99.9% identity
- lots of common variations, but also very rare ones
- variations ranging from 1 nucleotide (like SNPs) to tens of thousands nucleotides (structural variants)
- entire genes missing in some individuals (non-essential genes)

The "race" debate





© Daniel Utter

Other large scale human sequencing projects

- UK Biobank : to decipher the genomes of 500,000 individuals
- Iceland's effort to study the genomes of its entire human population.
- President Barack Obama's Precision Medicine Initiative, a genomics study of 1 million americans.

There's more than humans

- I5K: 5000 arthropodes
- Genome 10K: at least one individual from each vertebrate genus
- Vertebrate Genomes Project : find and sequence at least one individual from each of the approximately 66,000 vertebrate species *A Digital Noah's Ark Genome Library of Species*
- announcement of the Earth BioGenome Project (EBP) in Feb. 2017 at BioGenomics- Global Biodiversity Genomics Conference, (Smithsonian National Museum of Natural History | Washington, D.C)

What made this sequencing projects possible

The evolution of sequencing costs

HGP \$3 billion

Illumina, early 2000: \$10,000

Now, below \$1000

Cost per Raw Megabase of DNA Sequence



What accelerated the trend



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Practically, what brought the price of a genome down

- Better and cheaper sequencing machines
- Short genomic sequences ("reads") are cheaper to produce
- Better computing machines
- **Better algorithms**



NGS means high sequencing capacity







GS FLX 454 (ROCHE)



HiSeq 2000

(ILLUMINA)



5500xl SOLiD (ABI)



GS Junior







Ion TORRENT



Bioinformatics

Sequencing DNA and producing numerous DNA "texts" generated an increasing need for bioinformaticians in order to use: algorithms, statistics, and other mathematical techniques

to decipher the language of DNA deciphering by comparison

The key is comparison

E.g. 2 insect genomes suspected to be somewhat related, evolutionarily speaking

- a fruit fly (Drosophila melanogaster) and
- a malaria mosquito (Anopheles gamibae)

We would like to know what parts of the fruit fly genomic sequence are dissimilar and what parts are similar to the mosquito genomic sequence ?





What do we do from here

• medicine

- personalized medicine medical decisions based on a patient's predicted response or risk of disease
- therapeutic purposes study of pathogenic genomes like variants of E. coli, the antibiotics resistant strains of S. aureus, ...
- agriculture
 - identification and manipulation of genes linked to specific phenotypic traits, breeding by marker-assisted selection of variants, ...
- understanding microbial communities through metagenomics with impact on earth sciences, biomedicine, bioenergy, biotechnology, ...



How do we get from "reads" to complete sequences ?

Assembly Problem



Assembly is an extremely hard problem

Genomes cannot be read as reading a book, they are read by small pieces.

- Though becoming longer nowadays, reads remain quite short w.r.t. genome lengths
- There are lots of pieces !!!
- Assembly is not like solving a jigsaw puzzle
 - reads may have errors (up to 3%)
 - **repeats** are frequent (more than 50% of the human genome)
 - reads overlap
 - pieces from the genome may be missing
 - the reference may be different from the actual sequence or
 - the reference may not even exist

Assembly is an extremely hard problem

Not a jigsaw puzzle, more like a triazzle





Assembly problem

Input : a set of reads that are sub-strings of the genome

Output : the genome sequence explaining all the reads

Toy example :

Input : {GAAG, AAGT, GTAG, TAGA} Output : GAAGTAGA

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Assembly problem

Toy example :

Input : {GAAG, AAGT, GTAG, TAGA} Output : GAAGTAGA

AAGTAGAGTAGAAG is also a solution

Which one is the best?

Shortest does not necessarily mean best but fixing another criterion is difficult.

Shortest Common Superstring (SCS) problem



Computing a SCS solution as a Hamiltonian path

Equivalent to finding a Hamiltonian path in an overlap graph : visit every node in the graph exactly once

NP-hard problem

Approximation algorithms exist with factors : 4, 3, 2.89, 2.75, 2.5, 2.366, ...

The greedy method (4 approximation) :

- finds pairs of strings that overlap the best
- merges them
- repeats the operation

A particular case : r-SCS problem



SCS problem applied on a set of strings having the same length r.

Overlap Graph built on a set SE





Superstring solution

|| CGCAG || = 5



Superstring solution

|| CGCAGTC || = 7



Superstring solution

|| CGCAGTCAGCA || = 11



Superstring solution

CGCAGTCAGCATAA	= 14
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Superstring solution

|| CGCAGTCAGCATAACGCA || = 18



Superstring solution

||CGCAGTCAGCATAACGCATG|| = 20



Superstring solution

||CGCAGTCAGCATAACGCATGCGCAT|| = 25



Superstring solution

||CGCAGTCAGCATAACGCATGCG CAT|| = 25

max-length superstring solution (35)hamiltonian path length (10)

SCS solution = heaviest Hamiltonian path



Shortest superstring solution

||CGCAGCAGTCATAACGCATG|| = 20

= max-length superstring solution (35) hamiltonian path length (15)

MAX - ATSP problem

On the approximation ratio of the r-SCS problem
r-SCS problem : Golovnev et al. solution



Simple idea : focus on the best overlaps (r-1 overlaps)

- Build the de Brujn graph dB(r 1) on S and transpose the r-SCS problem into a 2-SCS instance.
- 2. Solve the 2-SCS instance with the algorithm from Crochemore *et al*.
- 3. Output the corresponding superstring solution for the original r-SCS problem.

de Bruijn Graph for k = 3



de Bruijn Graph

- overlap graph built on all the k-length strings present in the reads
- Nodes = k-mers and edges = (k-1)-overlaps

r-SCS problem : Golovnev et al.

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Build a de Bruijn Graph



r-SCS to 2-SCS



From r-SCS to 2-SCS



r-SCS problem on S : Golovnev et al.

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Solve 2-SCS : eulerian path with minimal additional edges



r-SCS problem on S : Golovnev et al.

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From 2-SCS back to r-SCS



Golovnev et al. - approximation ratio

$$\begin{array}{ll} \text{With} \quad x = \frac{w(H)}{n} & \text{we get the following ratios :} \\ \text{2-SCS based method :} & \frac{(r^2 - 2r + 2) - (r - 1)x}{r - x} & \text{MAX-ATSP} \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & &$$

Better than the general best SCS ratio (2.366) for r < 7

r-SCS problem : a hierarchical solution



Extend the idea of Golovnev et al. :

the best overlaps (r-1 overlaps) + the second best (r-2 overlaps)

r-SCS problem : a hierarchical solution

Step 1

Step 2

 Build the de Brujn graph dB(r – 1) on S and transpose the r-SCS problem into a 2-SCS instance.

Solve the 2-SCS instance with the algorithm from Crochemore *et al*.

Build a set of contigs S' by removing the edges added by the eulerian procedure; build a de Bruijn graph dB(r - 2) on the (r-2)-affixes of S' and transpose it into a 2-SCS instance.

Solve the novel 2-SCS instance and output the corresponding superstring solution (named γ) for the original r-SCS problem.

r-SCS hierarchical solution : Step 1 (Golovnev et al.)



r-SCS problem : a hierarchical solution

Step

4.

 Build the de Brujn graph dB(r – 1) on S and transpose the r-SCS problem into a 2-SCS instance.

2. Solve the 2-SCS instance with the algorithm from Crochemore *et al*.

Build a set of contigs S' by removing the edges added by the eulerian procedure; build a de Bruijn graph dB(r - 2) on the (r-2)-affixes of S' and transpose it into a 2-SCS instance.

Solve the novel 2-SCS instance and output the corresponding superstring solution (named γ) for the original r-SCS problem.

r-SCS hierarchical solution : Step 2 - build dB (r-2)



r-SCS problem : a hierarchical solution

Step

- Build the de Brujn graph dB(r 1) on S and transpose the r-SCS problem into a 2-SCS instance.
- 2. Solve the 2-SCS instance with the algorithm from Crochemore *et al*.
- 3. Build a set of contigs S' by removing the edges added by the eulerian procedure; build a de Bruijn graph dB(r 2) on the (r-2)-affixes of S' and transpose it into a 2-SCS instance.
- Solve the novel 2-SCS instance and output the corresponding superstring solution (named γ) for the original r-SCS problem.

r-SCS hierarchical solution : Step 2 - solve 2-SCS and compute γ



ACGCATGCGCAGCACAGTCCATAA

2-steps hierarchical solution - approximation ratio

With
$$x = \frac{w(H)}{n}$$
 we get the following ratios :

2-SCS based method
2-steps hierarchical method
$$\frac{(r^2 - 2r + 2) - (r - 1)x}{r - x}$$
MAX-ATSP
MAX-ATSP
MAX-ATSP
MAX-ATSP

$$\sum_{r=1}^{n} \left\{ \sum_{1 \le x \le r-1}^{n} \left\{ \min\left\{ \frac{(r^2 - 2r + 2) - (r - 1)x}{r - x}, \frac{r - \frac{2}{3}x}{r - x} \right\} \right\}$$

Better than the general best SCS ratio (2.366) for r < 7

SCS solution as an Eulerian path



No speech without the buzzword Big Data

Eulerian path is easier to compute but extremely large number of solutions and either way :

Graphs are long to build and extremely large to store hundreds of GB of data (hundreds of millions to billions of reads)

It's not BIG DATA, it's VERY BIG DATA

Clever solution for space-efficient de Bruijn graph representation

© Rizk & Chikhi 2013

Memory usage for k = 32, $N = 2 \times 10^9$, $E = 4 \times 10^9$

- Ascii sequence : 32B / node + 8B / edge (C pointer) 80 GB
- 2 bit / nt, 8 possible edges : (8 + 1)B / node [Z. Iqbal, 2012] **18 GB**

Edges can be inferred !

• Nodes only : 8B / node 16 GB

Self information of n nodes [Conway, Bromage, 2011]

• 20 bits per node 5 GB lower limit ???

Encoding de Bruijn graphs

Bloom filter

Bit array representing a set with a "precision" of ε . a proportion ε of elements will be wrongly included *false positives*

n elements : 1.44 $\log_2(1/\epsilon)^* n$ bits



Bloom filter

Hash function

- should have good "repartition" properties
- use of several functions to reduce false positive rate
 - insert n elements in a m bit array : ratio r = m/n
 - the FPR may be computed with respect to r



Bloom Filter

Set of nodes : {TAT, ATC, CGC, CTA, CCG, TCG, GCT}

de Bruijn graph as stored in a Bloom filter [*Pell et. al. 2012*]



Black nodes : true positives ; Red nodes : false positives

Bloom Filter

To traverse the graph from true positive nodes, only a small fraction of the false positives need to be avoided (critical FP).



Minia : De Bruijn representation by Chikhi & Rizk, 2013



Minia : rough performance comparison on human data



Minia : Besides genomic sequences assembly

Numerous applications and tools based on Minia :

- Metagenomic and transcriptome assembly
- Mutation detection
- Structural variant detection
- Targeted assembly
- K-mer counting
- Read error correction
- Read compression

Commet, Simka KisSplice, DiscoSNP TakeABreak, MindTheGap Mapsembler DSK Lordec, Bloocoo, LoRMA

Leon

Besides genomic sequences assembly



DiscoSNP



read	encoding			chor	dictionary
	anchor	left path	right path		0: AAAA 1: CGCT
CGCTAGATGA	1	o	6, A		2: CGAA 3: TAGG
GCTAGGTCTA	3	2	4, (T,3)		• • •

Minia : Besides genomic sequences assembly

Numerous applications and tools based on Minia :

GATB : the Genome Analysis Toolbox with de Bruijn graph read error correction

Read compression

Juce, DiscoSNP

TakeABreak, MindTheGap

Lordec, Bloocoo, LoRMA

Leon

Despite these clever ideas, things are still not simple

SCS - overly simplified abstraction of the assembly problem Remember the puzzle comparison

- reads may have errors thus overlaps may not be exact
- repeats are frequent (more than 50% of the human genome)
- reads may come from one strain or another

Things are not so simple

"Number of words consistent with genome graphs.

The size of the solution space for each chromosome using reads of length 50 nt. Only the 365 chromosomes with fewer than 2^{900} possible reconstructions are shown."

375 organisms (408 chromosomes in total)



© Kingsford, Schatz, Pop, 2010

Take home message

- Bioinformatics is undergoing a complete revolution
- Biological data and bioinformatics with it, are constantly evolving Tools and methods must be frequently updated
- Bioinformatics **does not mean** taking computer science methods and applying them blindly to biological data It means having truly understanding of biological data and producing methods that are perfectly adapted
- Large scale sequencing projects will shape the future of life related sciences


Thank you !

and

Happy Journées Montoises !