# RIGIDITY FOR SOME DYNAMICAL SYSTEMS OF ARITHMETIC ORIGIN

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## THE QUESTION OF RIGIDITY

Rejidity = there exists a sequence  $q_n \rightarrow \infty$  such that for any measurable set

 $\mu(T^{q_n}A\Delta A)\to \mathbf{0}.$ 

Veech (1982) : almost all interval exchanges are rigid.

Examples of <u>non-rigid</u> iet were known only for <u>3</u> intervals. Until Robertson (2017) and the square-tiled interval exchanges of F-H. (2016-17)

#### A THREE-INTERVAL EXCHANGE

Take the rotation  $Rx = x + \alpha$  modulo 1 and mark a point  $\beta$ .



#### THE EXAMPLE OF VEECH 1969

 $T(x,s) = (Rx, \sigma_0 s) \text{ if } x \text{ is in the interval } [0, \beta[\times\{s\}, \text{ assimilated with } [s-1, s-1+\beta[, T(x,s) = (Rx, \sigma_1 s) \text{ if } x \text{ is in the interval } [\beta, 1[\times\{s\}, \text{ assimilated with } [s-1+\beta, s[, T(x,s) = (Rx, \sigma_1 s) \text{ if } x \text{ is in the interval } [\beta, 1[\times\{s\}, \text{ assimilated with } [s-1+\beta, s[, T(x,s) = (Rx, \sigma_1 s) \text{ if } x \text{ is in the interval } [\beta, 1[\times\{s\}, \text{ assimilated with } [s-1+\beta, s[, T(x,s) = (Rx, \sigma_1 s) \text{ if } x \text{ is in the interval } [\beta, 1[\times\{s\}, 1[x], s]) \text{ if } x \text{ is in the interval } [\beta, 1[x], s] \text{ is in the interval } [\beta, 1[x],$ 



Thus the image intervals are

### **GENERALIZED VEECH**

We start from R, mark several points  $\beta_i$ , use permutations on  $\{1, \dots, d\}$ , take d copies of the interval [0, 1[. Optionally, change permutations at  $1 - \alpha$ , like in square-tiled iet.



#### **GRAND UNIFICATION**

We take  $\alpha$  irrational,  $0 = \beta_0 < \beta_1 < ... \beta_t < 1 - \alpha < \beta_{t+1} < ... \beta_r < \beta_{r+1} = 1$ ,  $\sigma_0$ , ...,  $\sigma_r$ ,  $\tau$ , permutations of  $\{1, ... d\}$ .

 $Rx = x + \alpha$  modulo 1.

$$T(x,s) = (x,\sigma_j s) \text{ if } \beta_j \leqslant x < \beta_{j+1}, \ j \neq t,$$
  

$$T(x,s) = (x,\sigma_t s) \text{ if } \beta_t \leqslant x < 1 - \alpha,$$
  

$$T(x,s) = (x,\tau s) \text{ if } 1 - \alpha \leqslant x < \beta_{t+1}.$$

Non-triviality conditions :  $\sigma_j \neq \sigma_{j+1}$ ,  $0 \leq j \leq r-1$ ,  $j \neq t$ ;  $\tau \neq \sigma_{t+1}$ ;  $\sigma_t \sigma_r \neq \tau \sigma_0$ .

#### SYMBOLIC SYSTEMS

Symbolic system = the shift on infinite sequences on a finite alphabet.

Trajectories =  $y_n = s_i$  if  $T^n y$  falls into the *i*-th interval in the *s*-th copy of [0, 1[.

A trajectory of T gives a trajectory of  $R: u \to \phi(u)$  by  $s_i \to i$ , for all i, s.

Linear recurrence of the coding = in the language of trajectories of R, every word of length n occurs in every word of length Kn.

## FIRST RESULT

**Theorem 1.** Under the non-triviality conditions and <u>minimality</u>, T is rigid (for any ergodic invariant measure) if  $\alpha$  has unbounded partial quotients, T is uniquely ergodic and non-rigid if the coding of R by the partition determined by  $\beta_1, ..., \beta_t, 1 - \alpha, \beta_{t+1}, ..., \beta_r$  is linearly recurrent.

Proofs as for square-tiled iet :

- an iet with LR behaves as a rotation with BPQ,
- when the non-triviality conditions are satisfied on every letter (ex : Veech 1969), T satisfies  $\overline{d}$ -separation,
- otherwise, average  $\overline{d}$ -separation.

#### MINIMALITY

For Veech 1969, it was known that if  $\beta$  is not in  $\mathbb{Z}(\alpha)$ , T is minimal. **Theorem 2.** T is minimal if and only if

$$1 \pm \beta = 2m\alpha + 2n$$

for some  $m \in \mathbb{Z}$ ,  $n \in \mathbb{Z}$ .

For the generalizations, we take all the  $\beta_i$  and  $\beta_i - \beta_j$  not in  $\mathbb{Z}(\alpha)$ . **Proposition 1.** An NCS for minimality is that no strict subset of  $\{1 \dots d\}$  is invariant by all the  $\sigma_i$  and  $\tau$ .

#### THE GREY ZONE AND OSTROWSKI

Case where  $\alpha$  has bounded partial quotients but the coding of R is <u>NOT</u> linearly recurrent.

Let  $\alpha$  be given,  $a_n$  its partial quotients : we use a form of alternating Ostrowski expansion of each  $\beta_i$  by  $\alpha$ , giving integers  $0 \leq b_n(\beta_i) \leq a_n$ . The Markov condition is  $b_n = a_n$ implies  $b_{n+1} = 0$ .

The integer  $b_{n+1}(\beta_i)$  tells us in which <u>column</u> is  $\beta_i$  for the *n*-th <u>Rokhlin tower</u> of the rotation *R*.

#### **ROKHLIN TOWERS**

*n* is fixed  $a = a_{n+1}$   $r = |q_n \alpha - p_n|$   $l = |q_{n-1} \alpha - p_{n-1}|$   $l' = |q_{n+1} \alpha - p_{n+1}|$ Then, up to the symmetry  $x \to -x$ , we have two towers.



#### BPQ BUT NOT LR

**Proposition 2.** For  $\alpha$  with BPQ, the coding of R is NOT LR if and only if

— either there exists *i* and pairs M, N with N-M arbitrarily large and  $b_m(\beta_i) = a_m - 1$ for  $M \leq m \leq N$ ,

 $eta_i$  is close to lpha

- or there exists *i* and pairs M, N with N-M arbitrarily large and  $(b_m(\beta_i), b_{m+1}(\beta_i)) = (a_m, 0)$  for each  $M \leq m = M + 2p \leq N$ ,  $\beta_i$  is close to  $\alpha$  (via 0)

- or there exist  $i \neq j$  and pairs M, N with N-M arbitrarily large and  $b_m(\beta_j) = b_m(\beta_i)$ for  $M \leq m \leq N$ .  $\underline{\beta_i}$  is close to  $\beta_j$ 

## IN THE GREY ZONE



#### NON-RIGID NON-LR

**Theorem 3.** If, whenever there is a run of m as in Proposition 1, either no  $\beta_i$  comes close to  $\alpha$ , or there exists  $\beta_j$  such that no  $\beta_k$ ,  $k \neq j$ , comes close to  $\beta_j$ , then T is not rigid.

This gives the first examples of not rigid not LR iet.

#### SO MANY SHADES OF GREY

**Theorem 4.** Suppose  $\sigma_k \sigma_j = \sigma_j \sigma_k$  for all  $j, k, \tau = \sigma_t$ . If infinitely often all the  $\beta_i$  come close to  $\alpha$  at the same time, then T is rigid.

This applies to Veech 1969, where d = 2: in the grey zone T is rigid.

**Proposition 3.** If there exist pairs M, N with N - M arbitrarily large, such that for all i either  $b_m(\beta_i) = a_m - 1$  for  $M \leq m \leq N$ , or  $(b_m(\beta_i), b_{m+1}(\beta_i)) = (a_m, 0)$  for each even  $M \leq m \leq N$ ,  $a_{N-1} \neq 1$  and  $a_N \neq 1$ , then T is rigid.

(Work in progress). When infinitely often all the  $\beta_i$  come close to  $\alpha$  at the same time, if Theorem 3 or 4 do not apply, and  $\alpha$  is the golden ratio, we can build rigid and non-rigid examples.

## RIGID LR

An interval exchange with a circular permutation is a rotation, thus is rigid, even if it is LR.

The Arnoux-Yoccoz interval exchange is semi-conjugate to a rotation of the 2- torus, which is rigid.

The measure-theoretic isomorphism between Arnoux-Yoccoz and the rotation, announced by Arnoux - Bernat - Bressaud (2011), was proved by Cassaigne (2018).

Thus Arnoux - Yoccoz is LR (self-induced, even) and rigid.