

On the group of a rational maximal bifix code

Alfredo Costa, Centre for Mathematics, University of Coimbra
joint work with Jorge Almeida, Revekka Kyriakoglou & Dominique Perrin

17e journées montoises d'informatique théorique — 2018



Centre for Mathematics
University of Coimbra

- Relativization of maximal bifix codes:

filtration

$$X = Z \cap F$$

of a maximal bifix code Z by a (uniformly) recurrent set F

- Relativization of descriptors of maximal bifix codes:

- $d(Z) \rightsquigarrow d_F(X)$
- $G(Z) \rightsquigarrow G_F(X)$

- We give necessary and sufficient conditions for

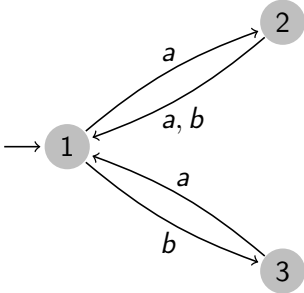
$$G(Z) \cong G_F(X)$$

- Methodological novelty: the use of free profinite monoids

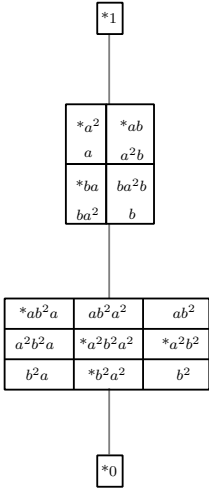
The syntactic monoid $M(L)$ of a language L is the transition monoid of the minimal automaton of L

We denote the map $A^* \rightarrow M(L)$ by η_L

e.g., $L = \{aa, ab, ba\}^*$



$M(L)$:



Green's relations

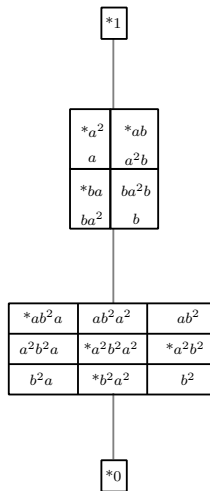
- 1 $u \mathcal{R} v \Leftrightarrow uM = vM$
- 2 $u \mathcal{L} v \Leftrightarrow Mu = Mv$
- 3 $u \mathcal{J} v \Leftrightarrow MuM = MvM$
- 4 $u \leq_{\mathcal{J}} v \Leftrightarrow MuM \subseteq MvM$
- 5 $\mathcal{H} = \mathcal{R} \cap \mathcal{L}$
- 6 In a finite monoid,

$$\mathcal{J} = \mathcal{R} \circ \mathcal{L} = \mathcal{L} \circ \mathcal{R} = \mathcal{R} \vee \mathcal{L}$$

All monoids we consider have this property.

- 7 For such monoids, if a \mathcal{J} -class contains an idempotent, then its \mathcal{H} -classes containing idempotents are isomorphic subgroups.

The abstract semigroup thus defined is the **Schützenberger group of the \mathcal{J} -class**.



Green's relations

- 1 $u \mathcal{R} v \Leftrightarrow uM = vM$
- 2 $u \mathcal{L} v \Leftrightarrow Mu = Mv$
- 3 $u \mathcal{J} v \Leftrightarrow MuM = MvM$
- 4 $u \leq_{\mathcal{J}} v \Leftrightarrow MuM \subseteq MvM$
- 5 $\mathcal{H} = \mathcal{R} \cap \mathcal{L}$
- 6 In a finite monoid,

$$\mathcal{J} = \mathcal{R} \circ \mathcal{L} = \mathcal{L} \circ \mathcal{R} = \mathcal{R} \vee \mathcal{L}$$

All monoids we consider have this property.

- 7 For such monoids, if a \mathcal{J} -class contains an idempotent, then its \mathcal{H} -classes containing idempotents are isomorphic subgroups.

The abstract semigroup thus defined is the **Schützenberger group of the \mathcal{J} -class**.

| | |
|---------------------------------|---------------------------------|
| *1 | |
| *a ² | *ab |
| ab ² a | a ² b ³ |
| a | a ² b |
| ab ² a ² | ab ² |
| a ² b ² a | ab ³ |
| ab ³ a | a ² b ² |
| *ba | ba ² b |
| b ³ a ² | *b ³ |
| ba ² | b |
| b ³ a | ba ² b ² |
| b ² a | b ² a ² b |
| b ² a ² | b ² |

Maximal bifix codes

A bifix code X of A^* is **maximal** if

$$X \subseteq Y \text{ and } Y \text{ is bifix} \implies X = Y$$

Motivation

- Perhaps the most studied and “tractable” class of codes.
- Important case: if $M(X^*)$ is a group, then X is a maximal bifix code, called a **group code**.

Filtration by recurrent sets

A bifix code X is **F -maximal** if $X \subseteq F$ and

$$X \subseteq Y \subseteq F \text{ and } Y \text{ is bifix} \implies X = Y$$

Theorem (Berstel & De Felice & Perrin & Reutenauer & Rindone; 2012)

If Z is maximal bifix, then $X = Z \cap F$ is...

- ... F -maximal bifix if F is recurrent
- ... finite if F is uniformly recurrent

If X is a rational maximal bifix code, then the degree of X is the rank $d(X)$ of the minimum ideal $J(X)$ of $M(X^*)$.

Theorem (The five authors of the 2012 paper + Dolce & Leroy; 2015)

If Z is a maximal bifix code and F is a tree set, then $X = Z \cap F$ is a basis of a subgroup of index $d(Z)$ of the free group $FG(A)$.

Let X be a rational maximal bifix code.

The F -minimum \mathcal{J} -class of $M(X^*)$, denoted $J_F(X)$, is the \mathcal{J} -class of $M(X^*)$ that is \mathcal{J} -minimum among the \mathcal{J} -classes intersecting $\eta_{X^*}(F)$.

The F -degree of X is the rank $d_F(X)$ of $J_F(X)$.

F -groups

The **group of Z** , denoted $G(Z)$, is the Schützenberger group of $J(Z)$.

The **F -group of X** , denoted $G_F(X)$, is the Schützenberger group of $J_F(X)$.

- $G(Z)$ is a permutation group of degree $d(Z)$
- $G_F(X)$ is a permutation group of degree $d_F(X)$

We want to relate $G(Z)$ with $G_F(X)$...

Theorem (Berstel & De Felice & Perrin & Reutenauer & Rindone; 2012)

If Z is a group code and F is Sturmian, then $G(Z) \simeq G_F(X)$ and $d(Z) = d_F(X)$.

(the theorem does not hold if we just replace “Sturmian” by “uniformly recurrent”, or “group code” by “maximal bifix code”)

F -groups

The **group of Z** , denoted $G(Z)$, is the Schützenberger group of $J(Z)$.

The **F -group of X** , denoted $G_F(X)$, is the Schützenberger group of $J_F(X)$.

- $G(Z)$ is a permutation group of degree $d(Z)$
- $G_F(X)$ is a permutation group of degree $d_F(X)$

We want to relate $G(Z)$ with $G_F(X)$...

Theorem (Berstel & De Felice & Perrin & Reutenauer & Rindone; 2012)

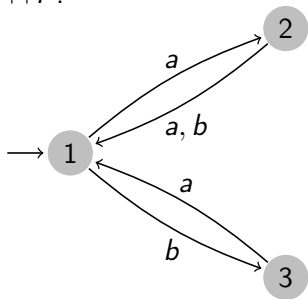
If Z is a group code and F is Sturmian, then $G(Z) \simeq G_F(X)$ and $d(Z) = d_F(X)$.

(the theorem does not hold if we just replace “Sturmian” by “uniformly recurrent”, or “group code” by “maximal bifix code”)

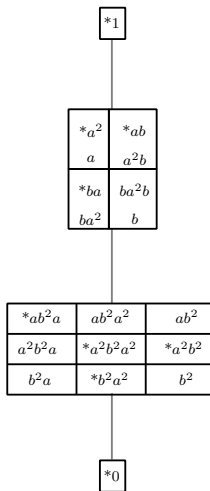
- $Z = \{aa, ab, ba, bb\}$
- $F = \text{“Fibonacci set”}$

$M(X^*)$:

Minimum automaton of X^* , where $X = Z \cap F$:



$$G(Z) \simeq G_F(X) \simeq \mathbb{Z}/2\mathbb{Z}$$



The profinite completion of A^*

If x, y are distinct elements of A^* , then there is a finite quotient $M = A^*/\sim$ separating x and y (i.e. $x \not\sim y$).

Let $r(x, y)$ be the smallest possible cardinal for M .

$$d(x, y) = 2^{-r(x, y)}$$

For this metric, let $\widehat{A^*}$ be the metric completion of A^* .

$\widehat{A^*}$ is a topological monoid, with a profinite topology.

More: $\widehat{A^*}$ is the **free profinite monoid generated by A** .

The group $G(F)$

If F is recurrent, then there is a \mathcal{J} -minimum \mathcal{J} -class contained in \overline{F} , denoted $J(F)$.

The group $G(F)$ is the (topological!) Schützenberger group of $J(F)$.

Tree and connected case

Theorem (Almeida & Costa; 2017)

If F is a tree set, then $G(F)$ is a free profinite group of rank $|A|$.

(the proof uses the "Return Theorem")

More precisely:

if F is a tree set, then $\pi|: G(F) \rightarrow \widehat{FG(A)}$ is an isomorphism:

$$\begin{array}{ccc} \widehat{A^*} & \longleftarrow & G(F) \\ & \searrow \pi & \downarrow \pi| \\ & & \widehat{FG(A)} \end{array}$$

More generally, $\pi|: G(F) \rightarrow \widehat{FG(A)}$ is onto if F is connected.

F -charged codes

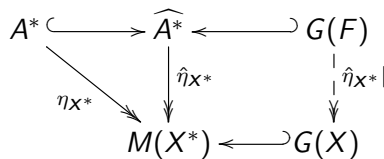
We say that a code X is:

- 1 F -charged if

$$\hat{\eta}_{X^*}(G(F)) = G(X)$$

- 2 weakly F -charged if

$$\hat{\eta}_{X^*}(G(F)) = G_F(X)$$



F -charged codes

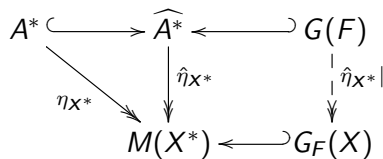
We say that a code X is:

- 1 F -charged if

$$\hat{\eta}_{X^*}(G(F)) = G(X)$$

- 2 weakly F -charged if

$$\hat{\eta}_{X^*}(G(F)) = G_F(X)$$



Necessary and sufficient conditions for $G(Z) \simeq G_F(Z \cap F)$

Theorem (Almeida & Costa & Kyriakoglou & Perrin; 2018)

Let:

- F recurrent
- Z rational maximal bifix code

Suppose also that $X = Z \cap F$ is rational.

The following conditions are equivalent:

- Z is F -charged
- $d_F(X) = d(Z)$, $G_F(X) \simeq G(Z)$ and X is weakly F -charged

Additionally: if F is *uniformly* recurrent, then the equality $d_F(X) = d(Z)$ is redundant in the second condition.

Necessary and sufficient conditions for $G(Z) \simeq G_F(Z \cap F)$

Theorem (Almeida & Costa & Kyriakoglou & Perrin; 2018)

Let:

- F recurrent
- Z rational maximal bifix code

Suppose also that $X = Z \cap F$ is rational.

The following conditions are equivalent:

- Z is F -charged
- $d_F(X) = d(Z)$, $G_F(X) \simeq G(Z)$ and X is weakly F -charged

Additionally: if F is **uniformly** recurrent, then the equality $d_F(X) = d(Z)$ is redundant in the second condition.

Group codes are “connected”-charged

Fact

If Z is a group code and F is connected, then Z is F -charged.

Proof.

$$\begin{array}{ccc} \widehat{A^*} & \xleftarrow{\quad} & G(F) \\ & \searrow \pi & \downarrow \pi| \\ & & \widehat{FG(A)} \\ \widehat{\eta}_{Z^*} \downarrow & & \\ M(Z^*) = G(Z) & & \end{array}$$

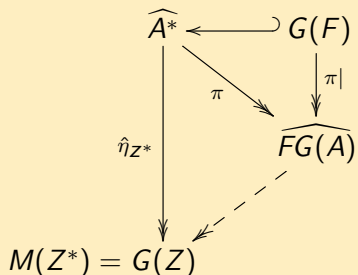


Group codes are “connected”-charged

Fact

If Z is a group code and F is connected, then Z is F -charged.

Proof.



Primitive substitutions

When $F = F_\varphi$ is described by a primitive substitution φ , we have an algorithm to decide if X is (weakly) F -charged: this is done via a profinite presentation for $G(F_\varphi)$, obtained by Almeida & Costa (2013).

Example

- F_τ : the Prouhet-Thue-Morse set, where

$$\tau : a \mapsto ab, b \mapsto ba$$

- Z : group code over $\{a, b\}$ generating the stabilizer of 1 via

$$a \mapsto (123), \quad b \mapsto (345)$$

- $G(Z) = A_5$
- Z is F_τ -charged

- $Z = \{aa, ab, ba\} \cup b^2(a^+b)^*b$

- $F = A^* \setminus A^*ab(b^2)^*aA^*$

Z is maximal bifix, but not a group code

F is recurrent, but not uniformly

Z is F-charged and $G(Z) \simeq G_F(X) \simeq S_3$

| | |
|---------------------------------|---------------------------------|
| *1 | |
| *a ² | *ab |
| ab ² a | a ² b ³ |
| a | a ² b |
| ab ² a ² | ab ² |
| a ² b ² a | ab ³ |
| ab ³ a | a ² b ² |
| *ba | ba ² b |
| b ³ a ² | *b ³ |
| ba ² | b |
| b ³ a | ba ² b ² |
| b ² a | b ² a ² b |
| b ² a ² | b ² |

| | | |
|----------------------------------|--|---|
| *b ⁴ a | b ² a ² b ² | b ⁴ a ² b |
| b ³ a ² | *a ³ | *b ³ ab |
| b ⁴ a ² | b ⁴ | b ⁴ ab |
| b ³ a | b ² a ² b ³ | b ³ a ² b |
| b ² a | b ³ a ² b ² | b ² a ² b |
| b ² a ² | b ² | b ² ab |
| *a ² | *a ⁴ | ab |
| ab ² a | a ² b ³ | a ² b |
| a | a ² b ⁴ | a ² b |
| ab ² a ² | ab ² | ab ² ab |
| a ² b ² a | ab ³ | ab ³ ab |
| ab ³ a | a ² b ² | a ² b ² ab |
| ba | ba ² b ⁴ | ba ² b |
| ba ² b ² a | *ba ³ b ² a | *ba ³ b ² ab |
| ba ² | ba ² b ⁴ | ba ² b |
| ba ³ a | ba ² b ² | ba ² b ² ab |
| ba ² a | ba ² b ³ | ba ² b ² a ² b |
| ba ² a ² | ba ² | ba ² b ² ab |

| | | | | | | | |
|------------------------------------|--------------------------------------|--|--|---|--|--|---|
| *a ² ba | *abab | a ² bab ³ | abab ² a | abab ² a ² b | abab ² a | abab ² | abab ² ab |
| *baaba | ba ² ab | *babab ² | ba ² b ² a | *babab ² ab | babab ² a | ba ² bab ² | ba ² bab ² ab |
| b ³ a ² ba | *b ³ abab | b ³ a ² bab ³ | b ³ a ² b ² a | b ³ abab ² a ² b | b ³ abab ² a | b ³ a ² bab ² | b ³ abab ² ab |
| ba ² b ² aba | *ba ² b ² abab | ba ² b ² abab ³ | ba ² b ² abab ² a | ba ² b ² abab ² a ² b | ba ² b ² abab ² a | ba ² b ² abab ² | ba ² b ² abab ² ab |
| ab ² aba | ab ² a ² bab | *ab ² abab ³ | *ab ² a ² bab ² a | ab ² abab ² ab | ab ² abab ² a ² | ab ² a ² bab ² | ab ² a ² bab ² ab |
| b ² aba | b ² a ² bab | b ² abab ² | b ² a ² bab ² a | b ² abab ² ab | *b ² abab ² a | b ² a ² bab ² | b ² a ² bab ² ab |
| a ² b ² aba | ab ² abab | a ² b ² abab ³ | a ² b ² abab ² a | a ² b ² abab ² a ² b | *a ² b ² abab ² a | *a ² b ² abab ² | a ² b ² abab ² ab |
| ba ² a ² ba | ba ² a ² bab | ba ² a ² bab ³ | ba ² a ² bab ² a | ba ² a ² bab ² a ² b | ba ² a ² bab ² a | ba ² a ² bab ² | ba ² a ² bab ² ab |

| | | | | | |
|---|---|---|---|---|--|
| *abab ² aba | abab ² a ² ba | abab ² abab | abab ² abab ² | abab ² abab ² a | abab ² abab ² ab |
| a ² ba ² aba | *abab ² aba | *a ² ba ² abab | a ² ba ² abab ² | a ² ba ² abab ² a | a ² ba ² abab ² ab |
| ba ² ab ² aba | *ba ² ab ² a ² ba | ba ² ab ² abab | *ba ² ab ² abab ² | ba ² ab ² abab ² a | *ba ² ab ² abab ² ab |
| b ² a ² ba ² aba | b ² a ² ba ² a ² ba | *b ² a ² ba ² abab | b ² a ² ba ² abab ² | b ² a ² ba ² abab ² a | b ² a ² ba ² abab ² ab |
| ab ² aba ² ba | ab ² aba ² a ² ba | ab ² aba ² abab | *ab ² aba ² abab ² | ab ² aba ² abab ² a | ab ² aba ² abab ² ab |

*0

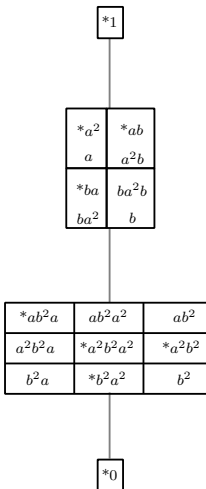
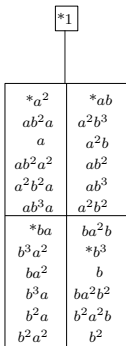
- $Z = \{aa, ab, ba\} \cup b^2(a^+b)^*b$

- $F =$ “Fibonacci set”

Z is maximal bifix, but not a group code

F is Sturmian

Z is not F -charged and $G(Z) \not\cong G_F(X)$



Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon“

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)

Conclusion: profinite is good!

We gave the first relevant “external” applications of the profinite Schützenberger group $G(F)$ of a uniformly recurrent set.

- the statement

$$G(Z) \cong G_F(Z \cap F) \text{ if } F \text{ is connected and } Z \text{ is group code}$$

uses no “profinite jargon”

- the definition of “ F -charged” gives a comprehensive framework to improve the latter
- the profinite group $G(F)$ serves as a sort of universal cover for F -groups
- “profinite stuff” facilitates synthetic statements
- advantages of the profinite monoid for proofs:
 - more “conceptual” proofs (diagram style)
 - enhanced combinatorics (“pseudowords” can be idempotent!)